

Fast computation of normal forms of polynomial matrices

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Polynomial matrix computations

Matrices over $\mathbb{K}[X]$
matrix $m \times m$

$$\begin{bmatrix} 3X + 4 & X^3 + 4X + 1 & 4X^2 + 3 \\ 5 & 5X^2 + 3X + 1 & 5X + 3 \\ 3X^3 + X^2 + 5X + 3 & 6X + 5 & 2X + 1 \end{bmatrix}$$

Fundamental operations

- multiplication
- kernel basis
- approximant basis

Transformation to normal forms

- triangularization \rightsquigarrow Hermite
- row reduction \rightsquigarrow Popov
- diagonalization \rightsquigarrow Smith

Polynomial matrix computations

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matrix $m \times m$

degree $d \rightsquigarrow \tilde{\mathcal{O}}(m^\omega d)$

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Matrices over $\mathbb{K}[X]$

matrix $m \times m$

degree $d \rightsquigarrow \tilde{O}(m^\omega d)$

type of average degree D/m

$$\begin{bmatrix} 3X + 4 & X^3 + 4X + 1 & 4X^2 + 3 \\ 5 & 5X^2 + 3X + 1 & 5X + 3 \\ 3X^3 + X^2 + 5X + 3 & 6X + 5 & 2X + 1 \end{bmatrix}$$

Fundamental operations

- multiplication
- kernel basis
- approximant basis

$\tilde{O}(m^\omega D/m)$ in specific cases

$\tilde{O}(m^\omega D/m)$

$\tilde{O}(m^\omega D/m)$

Transformation to normal forms

- triangularization \rightsquigarrow Hermite
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?

?

$\tilde{O}(m^\omega D/m)$

Hermite and Popov forms

working over $\mathbb{K} = \mathbb{Z}/7\mathbb{Z}$

$$\mathbf{A} = \begin{bmatrix} 3X + 4 & X^3 + 4X + 1 & 4X^2 + 3 \\ 5 & 5X^2 + 3X + 1 & 5X + 3 \\ 3X^3 + X^2 + 5X + 3 & 6X + 5 & 2X + 1 \end{bmatrix}$$

\rightsquigarrow using elementary row operations, transform \mathbf{A} into

Hermite form

$$\mathbf{H} = \begin{bmatrix} X^6 + 6X^4 + X^3 + X + 4 & 0 & 0 \\ 5X^5 + 5X^4 + 6X^3 + 2X^2 + 6X + 3 & X & 0 \\ 3X^4 + 5X^3 + 4X^2 + 6X + 1 & 5 & 1 \end{bmatrix}$$

Popov form

$$\mathbf{P} = \begin{bmatrix} X^3 + 5X^2 + 4X + 1 & 2X + 4 & 3X + 5 \\ 1 & X^2 + 2X + 3 & X + 2 \\ 3X + 2 & 4X & X^2 \end{bmatrix}$$

Example: constrained bivariate interpolation

As in Guruswami-Sudan list-decoding of Reed-Solomon codes

M of degree D ; L of degree $< D$

$$\mathbf{A} = \begin{bmatrix} M & & & & & \\ -L & 1 & & & & \\ -L^2 & & 1 & & & \\ \vdots & & & \ddots & & \\ -L^{m-1} & & & & & 1 \end{bmatrix}$$

Problem: find $\mathbf{p} = [p_1 \ \cdots \ p_m] \in \text{RowSpace}(\mathbf{A})$ such that

$$(\star) \quad \deg(p_j) < N_j \quad \text{for all } j$$

Approach:

- compute the Popov form \mathbf{P} of \mathbf{A} with degree weights on the columns
- return row of \mathbf{P} which satisfies (\star)

Shifted Popov form

Connects Popov and Hermite forms

$$\mathbf{s} = (0, 0, 0, 0) \quad \begin{bmatrix} [4] & [3] & [3] & [3] \\ [3] & [4] & [3] & [3] \\ [3] & [3] & [4] & [3] \\ [3] & [3] & [3] & [4] \end{bmatrix} \quad \begin{bmatrix} [7] & [0] & [1] & [5] \\ [0] & [1] & & [0] \\ & & [2] & \\ [6] & [0] & [1] & [6] \end{bmatrix}$$

Popov

$$\mathbf{s} = (0, 2, 4, 6) \quad \begin{bmatrix} [7] & [4] & [2] & [0] \\ [6] & [5] & [2] & [0] \\ [6] & [4] & [3] & [0] \\ [6] & [4] & [2] & [1] \end{bmatrix} \quad \begin{bmatrix} [8] & [5] & [1] & \\ [7] & [6] & [1] & \\ & & [2] & \\ [0] & [1] & & [0] \end{bmatrix}$$

s-Popov

$$\mathbf{s} = (0, D, 2D, 3D) \quad \begin{bmatrix} [16] & & & \\ [15] & [0] & & \\ [15] & & [0] & \\ [15] & & & [0] \end{bmatrix} \quad \begin{bmatrix} [4] & & & \\ [3] & [7] & & \\ [1] & [5] & [3] & \\ [3] & [6] & [1] & [2] \end{bmatrix}$$

Hermite

- normal form
- controlled average column degree
- and many useful properties

Shifted Popov form

For $\mathbf{A} \in \mathbb{K}[X]^{m \times m}$ nonsingular and $\mathbf{s} \in \mathbb{Z}^m$,
the **s-Popov form** of \mathbf{A} is the matrix $\mathbf{P} = \mathbf{UA}$ which is

$$\begin{array}{l} \mathbf{s}\text{-reduced} \\ \text{normalized} \end{array} \begin{bmatrix} [7] & [4] & [2] & [0] \\ [6] & [5] & [2] & [0] \\ [6] & [4] & [3] & [0] \\ [6] & [4] & [2] & [1] \end{bmatrix} \begin{bmatrix} [8] & [5] & [1] \\ [7] & [6] & [1] \\ & & [2] \\ [0] & [1] & & [0] \end{bmatrix}$$

sum of **diagonal degrees**:

$$d_1 + \cdots + d_m = \deg(\det(\mathbf{P})) = \deg(\det(\mathbf{A})) \leq D$$

Problem and previous work

Input: $\mathbf{A} \in \mathbb{K}[X]^{m \times m}$ nonsingular; shift $\mathbf{s} \in \mathbb{Z}^m$

Output: the \mathbf{s} -Popov form of \mathbf{A}

Previous fast algorithms focus on [Hermite](#) and [Popov](#) forms

[Popov](#) form: $\tilde{O}(m^\omega d)$, deterministic

[Giorgi-Jeannerod-Villard '03] [Sarkar-Storjohann '11] [Gupta-Sarkar-Storjohann-Valeriotte '12]

[Hermite](#) form: $\tilde{O}(m^\omega d)$, Las Vegas randomized

[Gupta-Storjohann '11] [Gupta '11]

$$\begin{cases} p_1 f_{11} + \cdots + p_m f_{1m} = 0 \pmod{M_1} \\ \vdots \\ p_1 f_{n1} + \cdots + p_m f_{nm} = 0 \pmod{M_n} \end{cases}$$

Reconstruction from equations

High-order lifting [Storjohann, 2003]

Reduction of basis matrix

$\deg(\mathbf{P}) \leq d$

\mathbf{P} triangular

Popov form

shifted
Popov form

Hermite form

Outline

- 1 reduction to **average degree** $d \in \mathcal{O}(D/m)$
- 2 Hermite form in $\tilde{\mathcal{O}}(m^\omega D/m)$, **deterministic**
- 3 **s**-Popov form in $\tilde{\mathcal{O}}(m^\omega D/m)$, probabilistic

1. Reduce to average degree

Example of **partial linearization** on the **columns** [Gupta et al., 2012]

$$\begin{bmatrix} (18) \\ [17] & (7) \\ [17] & [6] & (37) \\ [17] & [6] & [36] & (2) \end{bmatrix} \xrightarrow{\text{avg.}=16} \begin{bmatrix} (1) & [16] \\ [0] & [16] & (7) \\ [0] & [16] & [6] & (3) & [16] & [16] \\ [0] & [16] & [6] & [2] & [16] & [16] & (2) \end{bmatrix}$$

Elementary rows are inserted:

$$\begin{bmatrix} (1) & [16] & & & & & & \\ & X^{17} & -1 & & & & & \\ [0] & [16] & (7) & & & & & \\ [0] & [16] & [6] & (3) & [16] & [16] & & \\ & & & & X^{17} & -1 & & \\ & & & & & X^{17} & -1 & \\ [0] & [16] & [6] & [2] & [16] & [16] & (2) & \end{bmatrix}$$

↪ **preserves** determinant, Smith form, inverse...

1. Reduce to average degree

Problem: given \mathbf{A} and \mathbf{s} , find \mathbf{P}

using no field operation, build

- $\mathcal{L}(\mathbf{A}) \in \mathbb{K}[X]^{\tilde{m} \times \tilde{m}}$
- $\mathcal{L}(\mathbf{s}) \in \mathbb{Z}^{\tilde{m}}$

such that

- $\tilde{m} \leq 3m$ and $\deg(\mathcal{L}(\mathbf{A})) \leq \lceil D/m \rceil$,
- \mathbf{P} = submatrix of $\mathcal{L}(\mathbf{s})$ -Popov form of $\mathcal{L}(\mathbf{A})$

uses partial linearization techniques from [Gupta et al., 2012]

The bound D can be taken as the generic determinant degree:

$$\max_{\pi \in \text{Perm}(\{1, \dots, m\})} \sum_{1 \leq i \leq m} \overline{\deg}(a_{i, \pi_i})$$

$\rightsquigarrow D/m \leq$ average row and column degrees

$$\begin{cases} p_1 f_{11} + \cdots + p_m f_{1m} = 0 \pmod{M_1} \\ \vdots \\ p_1 f_{n1} + \cdots + p_m f_{nm} = 0 \pmod{M_n} \end{cases}$$

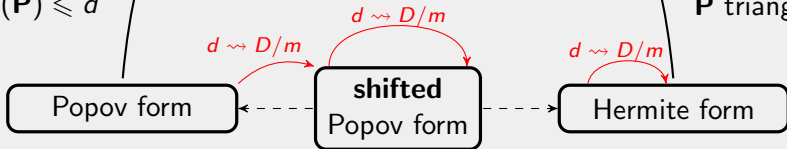
Reconstruction from equations

High-order lifting [Storjohann, 2003]

Reduction of basis matrix

$\deg(\mathbf{P}) \leq d$

\mathbf{P} triangular



2. Fast deterministic Hermite form

Previous fastest: $\tilde{O}(m^\omega d)$, Las Vegas

[Gupta-Storjohann, 2011]

Here: $\tilde{O}(m^\omega D/m)$, deterministic

(joint work with G. Labahn and W. Zhou [<http://arxiv.org/abs/1607.04176>])

Approach:

- 1 Find diagonal degrees [Zhou, 2012]
- 2 Reduce to Popov form computation

2.a. Find diagonal degrees

Partial computation of a triangularization:

$$\left[\begin{array}{c|c} \mathbf{A}_{11} & \mathbf{A}_{12} \\ \hline \mathbf{A}_{21} & \mathbf{A}_{22} \end{array} \right] \rightarrow \left[\begin{array}{c|c} \mathbf{B}_1 & \\ \hline * & \mathbf{B}_2 \end{array} \right] \rightarrow \left[\begin{array}{c|c|c} \mathbf{B}_{11} & & \\ \hline * & \mathbf{B}_{12} & \\ \hline & * & \mathbf{B}_{21} \\ \hline & & * & \mathbf{B}_{22} \end{array} \right] \rightarrow \dots$$

\rightsquigarrow yields diagonal entries in $\tilde{\mathcal{O}}(m^\omega d)$

- $\mathbf{B}_2 =$ small degree row basis of $\begin{bmatrix} \mathbf{A}_{12} \\ \mathbf{A}_{22} \end{bmatrix}$

[Zhou-Labahn, 2013]

- $\mathbf{N} =$ minimal kernel basis of $\begin{bmatrix} \mathbf{A}_{12} \\ \mathbf{A}_{22} \end{bmatrix}$

[Zhou-Labahn-Sorjohann, 2012]

- $\mathbf{B}_1 = \mathbf{N} \begin{bmatrix} \mathbf{A}_{11} \\ \mathbf{A}_{21} \end{bmatrix}$

2.b. Reduce to Popov form computation

$\mathbf{H} = -\mathbf{d}$ -Popov form of \mathbf{A} ($\mathbf{d} =$ diagonal degrees)

$$\mathbf{A} \xrightarrow{-\mathbf{d}\text{-reduction}} \mathbf{R} \xrightarrow[\text{(constant } \mathbf{U})]{\text{normalization}} \mathbf{H} = \mathbf{U}\mathbf{R}$$

$$\begin{bmatrix} [48] & [37] & [67] & [32] \\ [39] & [28] & [58] & [23] \\ [26] & [15] & [45] & [10] \\ [18] & [7] & [37] & [2] \end{bmatrix} \quad \begin{bmatrix} [18] & [7] & [37] & [2] \\ [18] & [7] & [37] & [2] \\ [18] & [7] & [37] & [2] \\ [18] & [7] & [37] & [2] \end{bmatrix} \quad \begin{bmatrix} (18) & & & \\ [17] & (7) & & \\ [17] & [6] & (37) & \\ [17] & [6] & [36] & (2) \end{bmatrix}$$

$-\mathbf{d}$ -reduction: via $\mathbf{0}$ -reduction \rightsquigarrow worst case $\tilde{O}(m^{\omega+1}d)$

normalization: in $\tilde{O}(m^{\omega}d)$

2.b. Reduce to Popov form computation

Partial linearization: (\mathbf{A}, \mathbf{d}) transformed into $(\mathcal{L}(\mathbf{A}), \mathcal{L}(\mathbf{d}))$

$$\left. \begin{array}{l} \mathcal{L}(\mathbf{A}) \text{ has degree } \leq d \\ \mathcal{L}(\mathbf{A}) \text{ has dimension } \leq 2m \\ \mathcal{L}(\mathbf{d}) \text{ has entries } \leq d \end{array} \right\} \Rightarrow \text{-}\mathcal{L}(\mathbf{d})\text{-reduction of } \mathcal{L}(\mathbf{A}) \text{ in } \tilde{\mathcal{O}}(m^\omega d)$$

$$\begin{array}{ccccc} \mathbf{A} & \xrightarrow{\text{-}\mathbf{d}\text{-reduction}} & \mathbf{R} & \xrightarrow[\text{(constant } \mathbf{U})]{\text{normalization}} & \mathbf{H} = \mathbf{UR} \\ \downarrow \text{partial linearization} & & & & \downarrow \text{partial linearization} \\ \mathcal{L}(\mathbf{A}) & \xrightarrow{\text{-}\mathcal{L}(\mathbf{d})\text{-reduction}} & \hat{\mathbf{R}} & \xrightarrow[\text{(constant } \hat{\mathbf{U}})]{\text{normalization}} & \mathcal{L}(\mathbf{H}) = \hat{\mathbf{U}} \hat{\mathbf{R}} \end{array}$$

\mathbf{H} directly obtained from $\mathcal{L}(\mathbf{H})$

$$\begin{cases} p_1 f_{11} + \cdots + p_m f_{1m} = 0 \pmod{M_1} \\ \vdots \\ p_1 f_{n1} + \cdots + p_m f_{nm} = 0 \pmod{M_n} \end{cases}$$

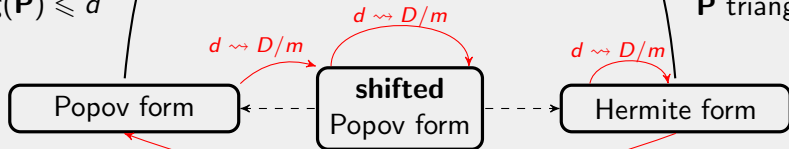
Reconstruction from equations

High-order lifting [Storjohann, 2003]

Reduction of basis matrix

$\deg(\mathbf{P}) \leq d$

\mathbf{P} triangular



via diagonal degrees

3. Fast \mathbf{s} -Popov form for arbitrary \mathbf{s}

Previous fastest: $\tilde{O}(m^\omega(d + \text{amp}(\mathbf{s}))) \subseteq \tilde{O}(m^{\omega+2}d)$,
relying on **non-shifted** Popov form computation [Gupta et al., 2012]

Here: $\tilde{O}(m^\omega D/m)$, Las Vegas randomized

Approach:

- a Build system of modular equations [Gupta-Storjohann, 2011]
- b Find \mathbf{s} -Popov basis of solutions [Neiger, 2016]

Note: yields fastest known algorithm for **Popov form** ($\mathbf{s} = \mathbf{0}$)

$$\begin{cases} p_1 f_{11} + \dots + p_m f_{1m} = 0 \text{ mod } M_1 \\ \vdots \\ p_1 f_{n1} + \dots + p_m f_{nm} = 0 \text{ mod } M_n \end{cases}$$

Reconstruction from equations

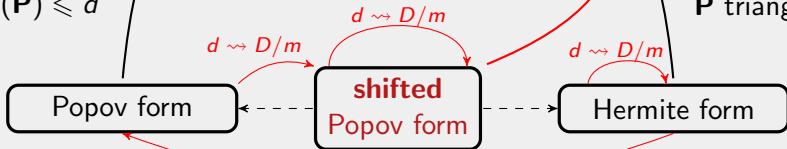
High-order lifting [Storjohann, 2003]

Reduction of basis matrix

Smith form of **A**
and reduced right
transformation

$\deg(\mathbf{P}) \leq d$

P triangular



via diagonal degrees

3.a. Build system of linear modular equations

Compute:

Smith form $\mathbf{UAV} = \text{diag}(1, \dots, 1, M_1, \dots, M_n)$

reduced right transformation $[\mathbf{0} \mid \mathbf{F}] = \mathbf{V} \bmod (1, \dots, 1, M_1, \dots, M_n)$

in probabilistic $\tilde{O}(m^\omega d)$ [Storjohann, 2003] [Gupta-Storjohann, 2011] [Gupta, 2011]

Then $\text{RowSpace}(\mathbf{A}) =$ all solutions $[p_1, \dots, p_m]$ to

$$\begin{cases} p_1 f_{11} + \dots + p_m f_{1m} = 0 \bmod M_1 \\ \vdots \\ p_1 f_{n1} + \dots + p_m f_{nm} = 0 \bmod M_n \end{cases}$$

\rightsquigarrow **s-Popov** form of $\mathbf{A} =$ **s-Popov** basis of solutions

3.b. Solve system of linear modular equations

Input: nonzero moduli M_1, \dots, M_n
system matrix $\mathbf{F} \in \mathbb{K}[X]^{m \times n}$
shift $\mathbf{s} \in \mathbb{Z}^m$

Output: the \mathbf{s} -Popov basis of $\{\mathbf{p} \mid \mathbf{p}\mathbf{F} = 0 \pmod{(M_1, \dots, M_n)}\}$

Result: $\tilde{\mathcal{O}}(m^\omega D/m)$ for arbitrary moduli, $n \in \mathcal{O}(m)$

where $D = \deg(M_1) + \dots + \deg(M_n)$

Previous work: $\tilde{\mathcal{O}}(m^\omega D/m)$ for

- Approximant bases: moduli = powers of X
- Interpolant bases: moduli given by roots and multiplicities
- Single degree-constrained solution (via structured system solving)

3.b. Solve system of linear modular equations

divide-and-conquer on the number of equations using ideas from

- [Jeannerod et al., 2016] (manage arbitrary shifts)
- [Gupta-Storjohann, 2011] (solution when diagonal degrees are known)

↪ remains the base case: one equation

$$p_1 f_1 + \cdots + p_m f_m = 0 \pmod{M}$$

P the sought s-Popov solution basis:

$$\mathbf{PF} = \begin{bmatrix} q_1 \\ \vdots \\ q_m \end{bmatrix} M \quad \Leftrightarrow \quad [\mathbf{P} \quad \mathbf{q}] \begin{bmatrix} \mathbf{F} \\ M \end{bmatrix} = 0$$

3.b. Solve system of linear modular equations

Reduction to **approximant basis**:

$$\begin{bmatrix} \mathbf{P} & \mathbf{q} \\ * & * \end{bmatrix} \begin{bmatrix} \mathbf{F} \\ M \end{bmatrix} = 0 \pmod{X^{\text{amp}(\mathbf{s})+2D}}$$

where $\text{amp}(\mathbf{s}) = \max(\mathbf{s}) - \min(\mathbf{s})$

New **divide-and-conquer** approach:

Recursion: $\mathbf{s} = (\mathbf{s}^{(1)}, \mathbf{s}^{(2)})$, $\mathbf{F} = \begin{bmatrix} \mathbf{F}^{(1)} \\ \mathbf{F}^{(2)} \end{bmatrix}$ with $\text{amp}(\mathbf{s}^{(i)}) \approx \text{amp}(\mathbf{s})/2$

Base case: $\text{amp}(\mathbf{s}) \in \mathcal{O}(D)$, cost $\tilde{\mathcal{O}}(m^\omega D/m)$ [Jeannerod et al., 2016]

3.b. Solve system of linear modular equations

- ① recursive call to find **splitting index** and $\mathbf{P}^{(1)}$:

$$\begin{bmatrix} \mathbf{P}^{(1)} & \mathbf{0} \\ * & * \end{bmatrix} = \mathbf{s}^{(1)}\text{-Popov sol. basis for } (\mathbf{F}^{(1)}, M) \rightsquigarrow \text{UpdateSplit}(\mathbf{s}, \mathbf{F})$$

- ② residual computation thanks to **known** $\mathbf{P}^{(1)}$:

$$\mathbf{A} = \begin{bmatrix} \mathbf{P}^{(1)} & \mathbf{0} & \mathbf{q}^{(1)} \\ * & \mathbf{P}^{(0)} & * \\ * & \mathbf{0} & q \end{bmatrix} = \mathbf{u}\text{-Popov app. basis for } \begin{bmatrix} \mathbf{F}^{(1)} \\ \mathbf{F}^{(2)} \\ M \end{bmatrix} \rightsquigarrow \begin{bmatrix} \mathbf{0} \\ \mathbf{G} \\ N \end{bmatrix} = \mathbf{A} \begin{bmatrix} \mathbf{F}^{(1)} \\ \mathbf{F}^{(2)} \\ M \end{bmatrix}$$

- ③ recursive call to find $\mathbf{P}^{(2)}$

$$\mathbf{P}^{(2)} = \mathbf{v}\text{-Popov sol. basis for } (\mathbf{G}, N), \text{ where } \text{amp}(\mathbf{v}) \approx \text{amp}(\mathbf{s})/2$$

- ④ compute $\mathbf{P} = \begin{bmatrix} \mathbf{P}^{(1)} & \mathbf{0} \\ * & \mathbf{P}^{(2)}\mathbf{P}^{(0)} \end{bmatrix}$ using **known diagonal degrees**

Conclusion

Linear systems of modular equations

- $\tilde{\mathcal{O}}(m^\omega D/m)$, deterministic ($n \in \mathcal{O}(m)$)
- return **s**-Popov solution basis for **arbitrary moduli**

Shifted row reduction of polynomial matrices

- $\tilde{\mathcal{O}}(m^\omega D/m)$, Las Vegas randomized
- computes **s**-Popov form for an **arbitrary shift**
- **Hermite** form: **deterministic**

Questions:

- **removing** the assumption $n \in \mathcal{O}(m)$?
- **deterministic** $\tilde{\mathcal{O}}(m^\omega D/m)$ Popov form?
- fast **deterministic** shifted Popov form?